

Ideas for a Hamiltonian Psychiatry

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0. Summary.

Relativity Theory, Freud's First Topic and the plastic neuronic nets could have a same Hamiltonian structure where Matte Blanco's symmetry principle and all the consequent bi-logic are particular cases of Freudian removal.

1. Nature 'Relativistic' and 'Freudian' of Plastic Neuronic Nets

Two mappings are given in Grappone [1998]. The former between the contents of Freud's first topic (mind's division in conscious, preconscious, unconscious, removal as operation of transferring mind objects from conscious into unconscious) and the contents of Einstein's relativity theory, the latter between the last contents and the that ones of the theory of the plastic neuron nets (i.e. of the neuron nets with variable connections). These two mappings give a direct mapping between contents of Freud's first topic and the contents of theory of the plastic neuron nets by composition. Precisely:

Freud's PsychoDynamics	Relativity Theory	Some Plastic Neuronic Nets
brain maximum speed : cleverness?	light speed	connection maximum speed
principle of the smallest anguish	Hamilton's principle	Hamilton's principle
<i>contradiction in unconscious</i>	gravitational attraction	engram fusion
<i>condensation</i>	gravitational attraction	engram fusion
<i>displacement</i>	gravitational attraction	engram connection
<i>internal instead of external reality</i>	relativistic mass to infinite	whole connection without inputs
<i>time absence in unconscious</i>	time dilatation	time dilatation
<i>space absence in unconscious</i>	Lorentz-Fitzgerald's contraction	Lorentz-Fitzgerald's contraction
anguish	momentum	connection momentum
irremovability	relativistic mass	neuron membership of an engram
processing speed	speed	connection speed
meaning	position vector	engram position vector
time	time	time

Table 1

Matte Blanco [1975] reduced all Freud's classical laws of unconscious (in italic in table 1) to the symmetry principle: *the unconscious makes commutative some non commutative relations with the consequence that it considers distinct mind objects (but in general quite similar) as a same mind object.* Now we would like to prove that Matte Blanco's symmetry principle (=unification of distinct terms) is a direct consequence of Freud's removal (i.e. to transfer objects from conscious into unconscious).

2. Standard Sentence Logic and Atomic Products

Let a, b, c, \dots be sentences. If α and β are sentences then $N\alpha, V\alpha\beta, A\alpha\beta, B\alpha\beta, C\alpha\beta, D\alpha\beta, E\alpha\beta, F\alpha\beta, G\alpha\beta, H\alpha\beta, I\alpha\beta, J\alpha\beta, K\alpha\beta, L\alpha\beta, M\alpha\beta, X\alpha\beta, O\alpha\beta$ are sentences in Polish language. they mean:

Polish language	Standard Language	Meaning
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$N\alpha$	$\sim \alpha$	the negation of α
$V\alpha\beta$	$(\alpha\vee\sim\alpha)\wedge(\beta\vee\sim\beta)$	the tautology on α and β
$A\alpha\beta$	$\alpha\vee\beta$	α and/or β i. e. the logical sum of α and β
$B\alpha\beta$	$\alpha\vee\sim\beta$	α is implied by β
$C\alpha\beta$	$\sim\alpha\vee\beta$	α implies β
$D\alpha\beta$	$\sim\alpha\vee\sim\beta$	the negation of α and/or the negation of β
$E\alpha\beta$	$(\alpha\wedge\beta)\vee(\sim\alpha\wedge\sim\beta)$	α is equivalent to β
$F\alpha\beta$	$\sim\alpha\vee(\beta\wedge\sim\beta)$	the negation of α
$G\alpha\beta$	$(\alpha\wedge\sim\alpha)\vee\sim\beta$	the negation of β
$H\alpha\beta$	$(\sim\alpha\vee\alpha)\wedge\beta$	β
$I\alpha\beta$	$\alpha\wedge(\sim\beta\vee\beta)$	α
$J\alpha\beta$	$(\sim\alpha\vee\sim\beta)\wedge(\alpha\vee\beta)$	α is equivalent to the negation of β
$K\alpha\beta$	$\alpha\wedge\beta$	α and β i. e. the logical products of α and β
$L\alpha\beta$	$\alpha\wedge\sim\beta$	α and the negation of β
$M\alpha\beta$	$\sim\alpha\wedge\beta$	the negation of α and β
$X\alpha\beta$	$\sim\alpha\wedge\sim\beta$	the negation of α and the negation of β
$O\alpha\beta$	$(\sim\alpha\wedge\alpha)\vee(\sim\beta\wedge\beta)$	the contradiction on α and β

Table 2

Call *atomic* either the sentences which are only represented by the small letters a, b, c, \dots , or their negations Na, Nb, Nc, \dots and Call *atomic product* any nidification of logical products of atomic sentences only $K\dots Ka\dots Nb\dots c\dots$. Call *disjunctive form* any nidification of logical sums of atomic products. All the dyadic connectives are equivalent to some disjunctive forms:

Na	Vab	Aab	Bab	Cab	Dab
Na	$AAAKabKaNb-KNabKNaNb$	$AAKabKaNb-KNabb$	$AAKabKaNb-KNaNb$	$AAKabKNab-KNaNb$	$AAKaNbKNab-KNaNb$
Eab	Fab	Gab	Hab	Iab	Jab
$AKabKNaNb$	$AKNabKNaNb$	$AKaNbKNaNb$	$AKabKaNb$	$AKabKaNb$	$AKaNbKNab$
Kab	Lab	Mab	Xab	Oab	
Kab	$KaNb$	$KNab$	$KNaNb$	$KaNa$	

Table 3

We observe quickly in table 3 that only Vab, Aab, Bab, Cab, Dab (see the first row couple) are dyadic connectives which can not be equivalent to disjunctive form with less than three atomic products.

Add to standard sentence logic the null sentence '\$' which is the neuter element of any dyadic connective, i. e. $N\$=\$, Va\$=V\$a=a, Aa\$=A\$a=a, Ba\$=B\$a=a, Ca\$=C\$a=a, Da\$=D\$a=a, Ea\$=E\$a=a, Fa\$=F\$a=a, Ga\$=G\$a=a, Ha\$=H\$a=a, Ia\$=I\$a=a, Ja\$=J\$a=a, Ka\$=K\$a=a, La\$=L\$a=a, Ma\$=M\$a=a, Xa\$=X\$a=a, Oa\$=O\$a=a$.

3. Standard Predicative Calculus of First Order: a not Standard Language

Let a, b, c, \dots be sentences when they are not in round brackets otherwise be terms. Let $(abc\dots)$ be atomic sentence when it is not in round brackets otherwise be term. If $\alpha, \beta, \gamma, \dots$ are terms then let $(\alpha\beta\gamma\dots)$ be atomic sentence when it is not in round brackets otherwise be term. If $(abc\dots)$ is sentence then let a be predicate otherwise be functional letter. Let (a) be always an individual constant. Let all the terms which are neither predicates nor functional letters nor individual constant be term variables. If a is a term variable and b is a sentence, then let $\prod ab$ and $\sum ab$ be sentences and mean respectively 'for every a, b is true' and 'for some a, b is true'. E. g. we have:

$$\begin{array}{l} \text{Our language:} \quad \prod x \sum y (axy) \quad \prod x \sum y K(ax)(by) \quad \prod x (ax(y)) \quad \prod x (ax(fzw)) \\ \text{Standard language:} \quad (\forall x)(\exists y) A_1^2(x, y) \quad (\forall x)(\exists y)(A_1^1(x) \wedge A_2^1(y)) \quad (\forall x) A_1^2(x, k) \quad (\forall x) A_1^2(x, \varphi_1^2(z, w)) \end{array}$$

4. Boolean relations between term variables

If a sentence a contains two distinct term variables b and c then if an atomic sentence of a contains either b or c then let Vbc be the Boolean relation between b and c in a ; otherwise eliminate any \prod or \sum with its first argument from a , after replace b to any atomic sentence in a with b , after replace c to any atomic sentence in a with c , after replace $\$$ to any atomic sentence in a without b and c , finally simplify a by standard sentence calculus and $\$$ rules (see the end of §2) till reducing a to a dyadic connective with b and c as arguments: let this connective be the Boolean relation between b and c in a .

E. g., consider the sentence $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$:

CALCULUS OF THE BOOLEAN RELATION BETWEEN m AND g IN $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$

$\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$	No atomic sentence contains either m or g :
$\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$	Elimination of any \prod or \sum with its first arguments:
$AAK(axm)N(axg)(apg)N(apm)$	Substitution of m to any atomic sentence with m :
$AAKmN(axg)(apg)Nm$	Substitution of g to any atomic sentence with g :
$AAKmNggNm$	No atomic sentence contains neither m nor g :
$AAKmNggNm$	Simplification by standard sentence calculus and $\$$ rules:
Vmg	

The Boolean relation between m and g in $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$ is Vmg .

CALCULUS OF THE BOOLEAN RELATION BETWEEN m AND p IN $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$

$\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(\underline{apm})$	An atomic sentence contains either m or p :
Vmp	

The Boolean relation between m and p in $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$ is Vmp .

CALCULUS OF THE BOOLEAN RELATION BETWEEN m AND x IN $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$

$\prod m \prod g \prod p \sum x AAK(\underline{axm})N(axg)(apg)N(apm)$	An atomic sentence contains either m or x :
Vmx	

The Boolean relation between m and x in $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$ is Vmx .

CALCULUS OF THE BOOLEAN RELATION BETWEEN g AND p IN $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$

$\prod m \prod g \prod p \sum x AAK(axm)N(axg)(\underline{apg})N(apm)$	An atomic sentence contains either g or p :
Vgp	

The Boolean relation between g and p in $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$ is Vgp .

CALCULUS OF THE BOOLEAN RELATION BETWEEN g AND x IN $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$

$\prod m \prod g \prod p \sum x AAK(axm)N(\underline{axg})(apg)N(apm)$	An atomic sentence contains either g or x :
Vgx	

The Boolean relation between g and x in $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$ is Vgx .

CALCULUS OF THE BOOLEAN RELATION BETWEEN p AND x IN $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$

$\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(\underline{apm})$	No atomic sentence contains either p or x :
$\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$	Elimination of any \prod or \sum with its first arguments:
$AAK(axm)N(axg)(apg)N(apm)$	Substitution of p to any atomic sentence with p :
$AAK(axm)N(axg)pNp$	Substitution of x to any atomic sentence with x :
$AAKxNxpNp$	No atomic sentence contains neither p nor x :
$AAKxNxpNp$	Simplification by standard sentence calculus and $\$$ rules:
Vpx	

The Boolean relation between p and x in $\prod m \prod g \prod p \sum x AAK(axm)N(axg)(apg)N(apm)$ is Vpx .

5. Variable Unification Metatheorem in Predicative Calculus of First Order

Variable unification metatheorem (VUM): GIVEN A SENTENCE $\prod b \prod c a(b,c)$ WITH a WHICH CONTAINS THE TERM VARIABLES b AND c , THE BOOLEAN RELATION BETWEEN b AND c IS A DYADIC CONNECTIVE WHICH IS EQUIVALENT TO SOME DISJUNCTIVE FORMS WITH LESS THAN THREE ATOMIC PRODUCTS IFF $\prod b a(b,b)$ AND $\prod c a(c,c)$ ARE THEOREMS IFF $\prod b \prod c a(b,c)$ IS THEOREM.

Proof:

Put $a(b,c) = Kd(b)e(c)$ with d which contains b but not c and e which contains c but not b , we have:

1a	$\prod b \prod c Kd(b)e(c)$	hypothesis
2a	$K \prod b d(b) \prod c e(b)$	1a, rules for a prenex normal form
3a	$\prod b Kd(b)e(b)$	2a, theorem: $E \prod b Kd(b)e(b) K \prod b d(b) \prod c e(b)$
4a	$C \prod b \prod c Kd(b)e(c) \prod b Kd(b)e(b)$	1a, 3a, deduction metatheorem
1b	$\prod b Kd(b)e(b)$	hypothesis

2b $K\lceil\lceil bd(b)\rceil\lceil be(b)\rceil$ 1b, theorem: $E\lceil\lceil bKd(b)e(b)K\lceil\lceil bd(b)\rceil\lceil be(b)\rceil$
 3b $\lceil\lceil b\lceil\lceil cKd(b)e(c)\rceil\rceil$ 2b, rules for a prenex normal form
 4b $C\lceil\lceil bKd(b)e(b)\rceil\lceil\lceil b\lceil\lceil cKd(b)e(c)\rceil\rceil$ 1b, 3b, deduction metatheorem

5ab $E\lceil\lceil b\lceil\lceil cKd(b)e(c)\rceil\lceil\lceil bKd(b)e(b)\rceil\rceil$ 4a, 4b, metatheorem: $Cab, Cba \not\vdash Eab$

Put $d(x) = Nf(x)$, we have:

6ab $E\lceil\lceil b\lceil\lceil cMf(b)e(c)\rceil\lceil\lceil bMf(b)e(b)\rceil\rceil$ 5ab, by substitution of $d(b)$ with $Nf(b)$

Put $d(b) = Nf(b)$ and $e(x) = Ng(x)$, we have:

7ab $E\lceil\lceil b\lceil\lceil cLd(b)g(c)\rceil\lceil\lceil bLd(b)g(b)\rceil\rceil$ 5ab, by substitution of $e(x)$ with $Ng(x)$

8ab $E\lceil\lceil b\lceil\lceil cXf(b)g(c)\rceil\lceil\lceil bXf(b)g(b)\rceil\rceil$ 7ab, by substitution of $d(b)$ with $Nf(b)$

Put $a(b,c) = Od(b)e(c)$ with d which contains b but not c and e which contains c but not b , we have:

1c $\lceil\lceil b\lceil\lceil cOd(b)e(c)\rceil\rceil$ hypothesis
 2c $\lceil\lceil bOd(b)e(b)\rceil$ 1c, Pseudoscotus' theorem: $CKaNab$
 3c $C\lceil\lceil b\lceil\lceil cOd(b)e(c)\rceil\lceil\lceil bOd(b)e(b)\rceil\rceil$ 1c, 2c, deduction metatheorem

1d $\lceil\lceil bOd(b)e(b)\rceil$ hypothesis
 2d $\lceil\lceil b\lceil\lceil cOd(b)e(c)\rceil\rceil$ 1d, Pseudoscotus' theorem: $CKaNab$
 3d $C\lceil\lceil bOd(b)e(b)\rceil\lceil\lceil b\lceil\lceil cOd(b)e(c)\rceil\rceil$ 1d, 2d, deduction metatheorem

4cd $C\lceil\lceil b\lceil\lceil cOd(b)e(c)\rceil\lceil\lceil bOd(b)e(b)\rceil\rceil$ 3c, 3d, metatheorem: $Cab, Cba \not\vdash Eab$

Put $a(b,c) = Ed(b)e(c)$ with d which contains b but not c and e which contains c but not b , we have:

1e $\lceil\lceil b\lceil\lceil cEd(b)e(c)\rceil\rceil$ hypothesis
 2e $\lceil\lceil cEd(b)e(c)\rceil$ 1e, Mendelson's axiom A4: $C\lceil\lceil ba(b)a(c)\rceil$
 3e $Ed(b)e(b)$ 2e, Mendelson's axiom A4: $C\lceil\lceil ba(b)a(c)\rceil$
 4e $\lceil\lceil bEd(b)e(b)\rceil$ 3e, generalization
 5e $C\lceil\lceil b\lceil\lceil cEd(b)e(c)\rceil\lceil\lceil bEd(b)e(b)\rceil\rceil$ 4e, deduction metatheorem
 1f $\lceil\lceil bEd(b)e(b)\rceil$ hypothesis
 2f $Ed(b)e(b)$ 1f, Mendelson's axiom A4: $C\lceil\lceil ba(b)a(c)\rceil$
 3f $NEd(b)Ne(b)$ 2f, tautology: $EEabNEaNb$
 4f $NEd(b)Ne(c)$ 3f, $CEd(b)e(c)Ed(b)e(b), ECabCNbNa$
 5f $Ed(b)e(c)$ 4f, tautology: $EEabNEaNb$
 6f $\lceil\lceil b\lceil\lceil cEd(b)e(c)\rceil\rceil$ 5f, generalization
 7f $\lceil\lceil bEd(b)e(b)\rceil \not\vdash \lceil\lceil b\lceil\lceil cEd(b)e(c)\rceil\rceil$ 6f, deduction methateorem ($\not\vdash$, not C , $\lceil\lceil c$ isn't in hypothesis)
 8ef $\lceil\lceil bEd(b)e(b)\rceil \leftrightarrow \lceil\lceil b\lceil\lceil cEd(b)e(c)\rceil\rceil$ 5e,7f, metatheorem: $(Cab \text{ or } (a \not\vdash b)), (Cba \text{ or } (b \not\vdash a)) \not\vdash (a \leftrightarrow b)$

Put $e(x) = Ng(x)$, we have:

9ef $\lceil\lceil bJd(b)g(b)\rceil \leftrightarrow \lceil\lceil b\lceil\lceil cJd(b)g(c)\rceil\rceil$ 8ef, tautology: $EEabJaNb$

Put $a(b,c) = Fd(b)e(c)$ with d which contains b but not c and e which contains c but not b , we have:

1g $\lceil\lceil b\lceil\lceil cFd(b)e(c)\rceil\rceil$ hypothesis
 2g $\lceil\lceil bNd(b)\rceil$ 1g, tautology: $EFabNa$
 3g $\lceil\lceil bFd(b)e(b)\rceil$ 2g, tautology: $EFabNa$
 4g $C\lceil\lceil b\lceil\lceil cFd(b)e(c)\rceil\lceil\lceil bFd(b)e(b)\rceil\rceil$ 1g, 4g, deduction metatheorem

1h $\lceil\lceil bFd(b)e(b)\rceil$ hypothesis
 2h $\lceil\lceil bNd(b)\rceil$ 1h, tautology: $EFabNa$
 3h $\lceil\lceil bFd(b)\rceil \lceil\lceil c e(c)\rceil$ 2h, tautology: $EFabNa$
 4h $\lceil\lceil b\lceil\lceil cFd(b)e(c)\rceil\rceil$ 3h, rules for a prenex normal form
 5h $C\lceil\lceil bFd(b)e(b)\rceil\lceil\lceil b\lceil\lceil cFd(b)e(c)\rceil\rceil$ 1h, 4h, deduction metatheorem

6gh $E\lceil\lceil bFd(b)e(b)\rceil\lceil\lceil b\lceil\lceil cFd(b)e(c)\rceil\rceil$ 4g, 5h, metatheorem: $Cab, Cba \not\vdash Eab$

Put $d(b) = Nf(b)$, we have:

7gh $E\lceil\lceil bIf(b)e(b)\rceil\lceil\lceil b\lceil\lceil cIf(b)e(c)\rceil\rceil$ 6gh, tautology: $EIabFNab$

Put $a(b,c) = Gd(b)e(c)$ with d which contains b but not c and e which contains c but not b , we have:

1i $\lceil\lceil b\lceil\lceil cGd(b)e(c)\rceil\rceil$ hypothesis
 2i $\lceil\lceil cNe(c)\rceil$ 1i, tautology: $EGabNb$
 3i $\lceil\lceil cGd(c)e(c)\rceil$ 2i, tautology: $EGabNb$
 4i $C\lceil\lceil b\lceil\lceil cGd(b)e(c)\rceil\lceil\lceil cGd(c)e(c)\rceil\rceil$ 1i, 3i, deduction metatheorem

1j	$\prod cGd(c)e(c)$	hypothesis
2j	$\prod cNe(c)$	1j, tautology: $EGabNb$
3j	$\prod cG\prod bd(b)e(c)$	2j, tautology: $EGabNb$
4j	$\prod b\prod cGd(b)e(c)$	3j, rules for a prenex normal form
5j	$C\prod cGd(c)e(c)\prod b\prod cGd(b)e(c)$	1j, 3j, deduction metatheorem
6ij	$E\prod b\prod cGd(b)e(c)\prod cGd(c)e(c)$	4i, 5j, metatheorem: $Cab, Cba \vee f Eab$

Put $e(x) = Ng(x)$, we have:

7ij	$E\prod b\prod cHd(b)g(c)\prod cHd(c)g(c)$	6ij, tautology: $EHabGaNb$
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k	$E\prod b\prod cC(abc)\prod cC(abb), E\prod b\prod cC(abc)\prod cC(acc)$	are not theorems: (let (abc) be 'b=c')
l	$E\prod b\prod cAd(b)e(c)\prod cAd(b)e(b), E\prod b\prod cAd(b)e(c)\prod cAd(c)e(c)$	are not theorems: (let d be Ne)
m	$E\prod b\prod cBd(b)e(c)\prod cBd(b)e(b), E\prod b\prod cBd(b)e(c)\prod cBd(c)e(c)$	are not theorems: (let d be e)
n	$E\prod b\prod cCd(b)e(c)\prod cCd(b)e(b), E\prod b\prod cCd(b)e(c)\prod cCd(c)e(c)$	are not theorems: (let d be e)
o	$E\prod b\prod cVd(b)e(c)\prod cVd(b)e(b), E\prod b\prod cVd(b)e(c)\prod cVd(c)e(c)$	are theorems (Vab is tautology), but two variables which have V as Boolean relation cannot be generally unified otherwise $CC\prod x(ax)\prod x(bx)\prod xC(ax)(bx)$ which is not theorem is proved.

* Consider $\prod b\prod c\Phi d(b)e(c)$ where Φ can be E or F or G or H or I or J or K or L or M or X or O and where $d(b)$ does not contain c and $e(c)$ does not contain b . b and c have always a pseudoboollean relation in $a(b,c)$ which is equivalent to a disjunctive form with minus than three atomic products (see §2). 5ab, 6ab, 7ab, 8ab, 4cd, 8ef, 9ef, 6gh, 7gh, 6ij, 7ij prove that $a(b,c)$ is theorem at less of substitions of b with c and vice versa. THUS VUM IS TRUE FOR $\prod b\prod c\Phi d(b)e(c)$.

** Let VUM be true for b and c in $a(b,c)$, d be sentence without b and c and Ψ be a dyadic connective at will. The Boolean relation of b and c is in $\Psi da(b,c)$: $\Psi da(b,c)$, i. e. $a(b,c)$, i. e. the same that in $a(b,c)$. Thus VUM is also true for $\Psi da(b,c)$ and, by iteration, VUM IS TRUE FOR ANY FORMULA WHICH CONTAINS $\prod b\prod c\Psi d(b)e(c)$ AND CONTAINS b AND c ONLY IN $\prod b\prod c\Psi d(b)e(c)$.

*** Any $a(b,c)$ where VUM is true for b and c has to be equivalent to an oportune $\Phi f(b)g(c)$ because b and c have to have a particular Boolean relation and not to stay in a same atomic sentence. So VUM IS TRUE FOR $\prod b\prod c\Phi f(b)g(c)$ WITH f AND g AT WILL.

**** The logical product od ** and *** is VUM.

5. Mind Contents in Freud's First Topic as Atomic products

A tautology has the complete disjunctive form because from it we can obtain the other sentences by elimination of one or more atomical products. Now put a mapping between standard sentence logic and Freud's first topic (conscious-preconscious-unconscious with removal and conscientization). Let any atomic product where any atomic sentence occurs one time at most be a mind content. Let the tautology be the mind situation with all the mind content in the conscious and the contradiction the mind situation with all the mind content in the unconscious. Thus all the other sentences correspond to mind situations with mind contents in the conscious and in the unconscious. So, if we represent Freud's first topic as

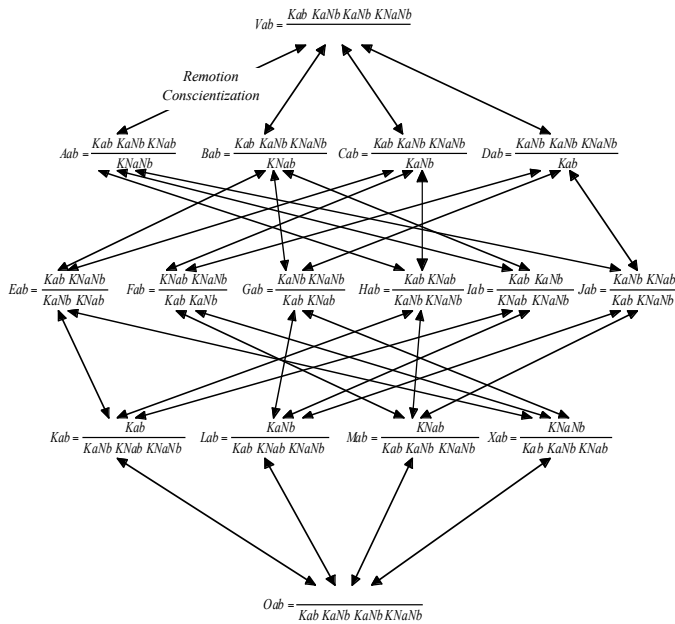
$preconscious \frac{conscious (conscient contents)}{unconscious (removed contents)}$, then we can rewrite the disjunct form of the dyadic tautology

Vab as $\frac{Kab \ KaNb \ KaNb \ KNaNb}{Kab \ KaNb \ KaNb \ KNaNb}$ (anything in conscious) and, consequently, the dyadic contradiction

Oab as $\frac{Kab \ KNaNb}{Kab \ KaNb \ KaNb \ KNaNb}$ (anything in unconscious). All the other sentences are intermediate.

E. g. Eab becomes $\frac{Kab \ KNaNb}{KaNb \ KbNa} \cdot \frac{Kab \ KNaNb}{KaNb \ KbNa} \xrightarrow{\text{remotion of } KNaNb} \frac{Kab}{KaNb \ KbNa \ KNaNb}$ is a mind

transition example by remotion. In this way Matte Blanco's symmetry becomes VUM. Here a simple dyadic model:



All possible atomic products in the conscious:
greatest thought freedom

Few disjuncts which are removed from conscious:
only few inhibitions and psychical limits

Atomic product removal is so big to permit the term
unification theorem, i. e. Matte Blanco's symmetry
principle, application:
neurotical processes

Not removed atomic product are so few that some ideas
cannot be changed and become delirium:
psychotical processes

Whole atomic product removal, i. e. contradiction, i. e.
conscious absence:
anything is anything for Pseudoscotus' theorem

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